

Interacting Frobenius Algebras and the Structure of Multipartite Entanglement

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Introduction

- ▶ For many quantum algorithms and protocols, entangled states are where the magic happens
- ▶ Often the study of entanglement is concerned with numeric measures, or exhaustive (and exhausting!) classification theorems
- ▶ While this is useful, we take a courser-grained approach, looking at *behavioural*, or *algebraic* properties of entangled states
- ▶ We show that a small number of interacting algebraic structures can describe a large class of multipartite entangled states
- ▶ The term “algebraic” is massively overloaded, so lets clarify

Entanglement and Algebra

- ▶ First note, we shall switch freely between the notion of a vector $|\psi\rangle \in H$ and its associated map

$$\psi : \mathbb{C} \rightarrow H :: 1 \mapsto |\psi\rangle$$

- ▶ Information flows *across* entangled states. Taking this seriously, we can regard states $|\Psi\rangle \in A \otimes B$ as maps

$$\tilde{\Psi} : A^* \rightarrow B :: \langle a| \mapsto (\langle a| \otimes 1)|\Psi\rangle$$

Entanglement and Algebra

- ▶ If we fix an isomorphism $\phi : A \cong A^*$, then $\tilde{\Psi}\phi$ is a new map:

$$\hat{\Psi} : A \rightarrow B$$

- ▶ Consider a tripartite state $|\Psi\rangle \in A \otimes A \otimes A$, then we can do this trick to get

$$\hat{\Psi} : A \otimes A \rightarrow A$$

- ▶ $\hat{\Psi}$ is a bilinear multiplication of vectors in A . If we impose conditions like associativity, etc. $(A, \hat{\Psi})$ becomes an *algebra* over A

LOCC and SLOCC

- ▶ We'll consider pure state entanglement on qubits
- ▶ States $|\Psi\rangle$ and $|\Phi\rangle$ are equivalent up to LOCC (local operations and classical communication) iff there exist unitaries $U_i : \mathbb{C}^2 \rightarrow \mathbb{C}^2$ such that

$$|\Psi\rangle = (U_1 \otimes \dots \otimes U_n) |\Phi\rangle$$

- ▶ $|\Psi\rangle$ and $|\Phi\rangle$ are equivalent up to stochastic LOCC, or SLOCC, if they can be interconverted, as above, but with non-zero probability
- ▶ This is equivalent to saying that for local *invertible* maps L_i

$$|\Psi\rangle = (L_1 \otimes \dots \otimes L_n) |\Phi\rangle$$

Entanglement, LOCC, and SLOCC

- ▶ There are two SLOCC classes in $\mathbb{C}^2 \otimes \mathbb{C}^2$, $|00\rangle$ and $|00\rangle + |11\rangle$
- ▶ In 3 qubits, there are only two genuine (non-product) states, up to SLOCC:

$$|\text{GHZ}\rangle = |000\rangle + |111\rangle \quad |\text{W}\rangle = |100\rangle + |010\rangle + |001\rangle$$

Associative Algebra

- ▶ $(A, *)$ is an *associative k -algebra* if A is a k -vector space and $(- * -) : A \times A \rightarrow A$ is an associative, bilinear multiplication defined on A
- ▶ It is *unital* if a vector $e \in A$ acts as a unit
- ▶ Since $*$ is bilinear, it factors uniquely through the tensor product

$$\left(A \times A \xrightarrow{*} A \right) = \left(A \times A \xrightarrow{\otimes} A \otimes A \xrightarrow{\mu} A \right)$$

- ▶ From now on, we shall focus on unital algebras on finite-dimensional, complex Hilbert spaces

Coalgebras

- ▶ So, we can alternately define a unital algebra $(H, *, |e\rangle)$ as a triple (H, μ, η) , where $\eta :: 1 \mapsto |e\rangle$ and
 - ▶ μ as associative: $\mu(1 \otimes \mu) = \mu(\mu \otimes 1)$
 - ▶ η is the unit: $\mu(1 \otimes \eta) = \mu(\eta \otimes 1) = 1$
- ▶ We could also define an algebra on the dual space $(H^*, \delta^*, \epsilon^*)$, or, equivalently, a *coalgebra* on H
- ▶ A (unital) *coalgebra* is a triple $(H, \delta : A \rightarrow A \otimes A, \epsilon : \mathbb{C} \rightarrow H)$ that is
 - ▶ *coassociative*: $(1 \otimes \delta)\delta = (\delta \otimes 1)\delta$
 - ▶ *counital*: $(\epsilon \otimes 1)\delta = (1 \otimes \epsilon)\delta = 1$

Associative Algebra (Graphically)

- ▶ An associative algebra is a triple

$$\left(H, \underbrace{\text{multiplication}}_{\text{multiplication}}, \underbrace{\text{unit}}_{\text{unit}} \right)$$

- ▶ Such that



And upside down...

- ▶ A coassociative coalgebra is a triple

$$\left(H, \underbrace{\begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} : H \rightarrow H \otimes H}_{\text{comultiplication}}, \underbrace{\begin{array}{c} \bullet \\ | \end{array} : H \rightarrow \mathbb{C}}_{\text{counit}} \right)$$

- ▶ Such that

The diagram shows two equations. The first equation is the coassociativity condition: a tree with a root node and three children, where the right child is itself a tree with a root and two children, is equal to a tree with a root and three children, where the left child is itself a tree with a root and two children. The second equation is the counit property: a tree with a root node and two children, both of which are themselves trees with a root and two children, is equal to a tree with a root node and two children, which is equal to a vertical line.

Commutative Frobenius Algebras, Definition 1

- ▶ Since we've fixed an isomorphism with the dual space, any algebra on H canonically yields a coalgebra. Equivalently, we could start with one and make some axioms
- ▶ A *commutative Frobenius algebra* is a quintuple:

$$(H, \mu, \eta, \delta, \epsilon)$$

- ▶ (H, μ, η) is a commutative algebra,
- ▶ (H, δ, ϵ) is a cocommutative coalgebra, and
- ▶ $(\mu \otimes 1)(1 \otimes \delta) = (1 \otimes \mu)(\delta \otimes 1) = \delta\mu$

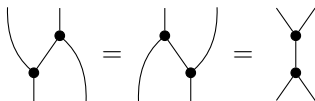
Commutative Frobenius Algebras, Definition 1

- ▶ A *commutative Frobenius algebra* is a quintuple:

$$(H, \begin{array}{c} \diagdown \\ \bullet \\ \diagup \end{array}, \begin{array}{c} \bullet \\ \uparrow \end{array}, \begin{array}{c} \bullet \\ \diagdown \\ \diagup \end{array}, \begin{array}{c} \downarrow \\ \bullet \end{array})$$

- ▶ Such that

- ▶ $(H, \begin{array}{c} \diagdown \\ \bullet \\ \diagup \end{array}, \begin{array}{c} \bullet \\ \uparrow \end{array})$ is a commutative algebra,
- ▶ $(H, \begin{array}{c} \bullet \\ \diagdown \\ \diagup \end{array}, \begin{array}{c} \downarrow \\ \bullet \end{array})$ is a cocommutative coalgebra, and

- ▶ 

Example

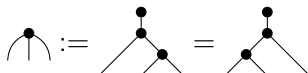
- ▶ For a finite-dimensional Hilbert space H , fix an ONB $\{|i\rangle\}$
- ▶ Let $\delta :: |i\rangle \mapsto |ii\rangle$, copy each of the basis vectors
- ▶ $\epsilon :: |i\rangle \mapsto 1$ uniformly deletes them
- ▶ $\mu :: |ii\rangle \mapsto |i\rangle, |ij\rangle \mapsto 0$ fuses them (aka, it performs the Schur product w.r.t. the basis)
- ▶ $\eta :: 1 \mapsto \sum |i\rangle$ is the unit of μ

Finding the Tripartite States

- ▶ We started by talking about tripartite states. What's that have to do with Frobenius algebras?
- ▶ A tripartite state $|\Psi\rangle \in H \otimes H \otimes H$ is a map

$$\Psi : \mathbb{C} \rightarrow H \otimes H \otimes H$$

- ▶ ...so is this:



- ▶ We can talk about the tripartite states *generated* by Frobenius algebras
- ▶ We're particularly interested in *symmetric* states, hence *commutative* Frobenius algebras

Commutative Frobenius Algebras, Definition 2

- ▶ A third (and final) definition...
- ▶ A *commutative Frobenius algebra* is a triple

$$(\Psi \in H \otimes H \otimes H, \phi : H \otimes H \rightarrow \mathbb{C}, \epsilon : H \rightarrow \mathbb{C})$$

- ▶ Such that
 - ▶ $(1 \otimes \phi \otimes 1)(\Psi \otimes \Psi)$ is symmetric, and
 - ▶ $\psi := (\epsilon \otimes 1 \otimes 1)\Psi$ generates an inverse to ϕ , or a “cap”. I.e.:


$$(1 \otimes \phi)(\psi \otimes 1) = 1$$

Commutative Frobenius Algebras, Definition 2

- ▶ A *commutative Frobenius algebra* is a triple

$$\left(\begin{array}{c} \bullet \\ \cap \\ | \\ | \\ | \end{array}, \begin{array}{c} \cup \\ \bullet \end{array}, \bullet \right)$$

- ▶ Such that

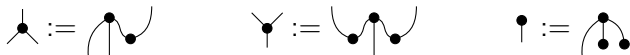
- ▶  is symmetric, and

- ▶  =

Equivalence of Definitions

► It is well known that (def 1) \iff (def 2)

► If we let:



then (def 2) \implies (def 1)

► The other way, let:



then (def 1) \implies (def 2)

Examples

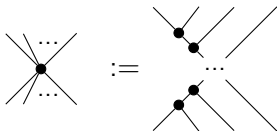
- ▶ Writing the Frobenius algebra as a triple $(\begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ \bullet \end{array}, \begin{array}{c} \diagdown \quad \diagup \\ \bullet \end{array}, \begin{array}{c} \bullet \\ \diagdown \\ \bullet \end{array})$, we see some familiar faces...

▶
$$\left(\underbrace{|000\rangle + |111\rangle}_{\text{GHZ state}}, \underbrace{\langle 00| + \langle 11|}_{\text{Bell state}}, \langle +| \right)$$

▶
$$\left(\underbrace{|100\rangle + |010\rangle + |001\rangle}_{\text{W state}}, \underbrace{\langle 10| + \langle 01|}_{\text{EPR state}}, \langle 0| \right)$$

Normal Forms

- ▶ These equivalent definitions are reflective of the fact that Frobenius algebras are largely topological in nature
- ▶ For “canonical trees,” we use the following notation:



- ▶ We call this thing a *spider*
- ▶ In fact, any connected, *acyclic* graph can be rewritten as such a tree. The equivalences on the previous slide are all special cases of this.
- ▶ ...and *any* graph can be rewritten as a tree with some loops in the middle

Special and Anti-special

- ▶ The only thing left to determine is what happens at the loops
- ▶ There's a lot of choices, but these two are particularly interesting:



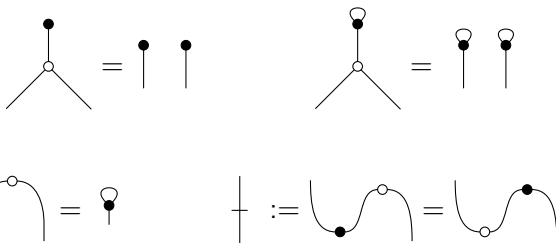
- ▶ If a CFA satisfies the left property, we call it *special*
- ▶ If it satisfies the right, we call it *anti-special*

The Punchline

- ▶ Why do we like these? Two classification results:
 - ▶ **Theorem:** A commutative Frobenius algebra over \mathbb{C}^2 is locally convertible to a special one *if and only if* its induced tripartite state is SLOCC with GHZ.
 - ▶ **Theorem:** A commutative Frobenius algebra over \mathbb{C}^2 is locally convertible to an anti-special one *if and only if* its induced tripartite state is SLOCC with W.
- ▶ Since all such states are genuine tripartite, and GHZ and W are the only tripartite qubit states, we have an easy corollary:
 - ▶ **Corollary:** Every CFA over \mathbb{C}^2 is either locally equivalent to a special CFA or to an anti-special CFA.

Interaction Rules

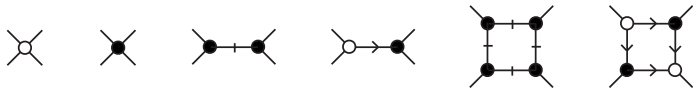
- ▶ Why should you care?
- ▶ If we denote the GHZ algebra by $(\text{Y}, \text{I}, \text{M}, \text{O})$ and the W algebra by $(\text{Y}, \text{I}, \text{M}, \text{O})$, we notice some interaction properties:



- ▶ And many more. Some we don't know yet!

Four-Partite and Beyond

- ▶ Looking at these interaction behaviours, we found representatives for lots of 4-partite SLOCC classes:



- ▶ As well as came up with an inductive technique for cooking them up...

Inductive Classification

- ▶ In 2006, L. Lamata, J. Leon, D. Salgado, E. Solano published a paper called *Inductive classification of multipartite entanglement under SLOCC*, where they identified an inductive technique for classifying multipartite systems
- ▶ They did this by regarding a state in $(\mathbb{C}^2)^{\otimes n}$ as a map from $(\mathbb{C}^2)^{\otimes(n-1)}$ to \mathbb{C}^2 , and looking at the vectors spanning the right singular subspace \mathfrak{W}
- ▶ Since these vectors are in $(\mathbb{C}^2)^{\otimes(n-1)}$, they apply the inductive procedure down to a base case of $n = 2$, where the SLOCC classes are trivially computable

Inductive Classification

- ▶ At $n \geq 4$, there are infinitely many SLOCC classes, so Lamata and others introduce free parameters into the classification scheme, identifying SLOCC “superclasses”
- ▶ The representatives of these n partite classes are all of the form $|0, \Psi_0\rangle + |1, \Psi_1\rangle$. So, to represent this graphically, we would like to have some kind of map:

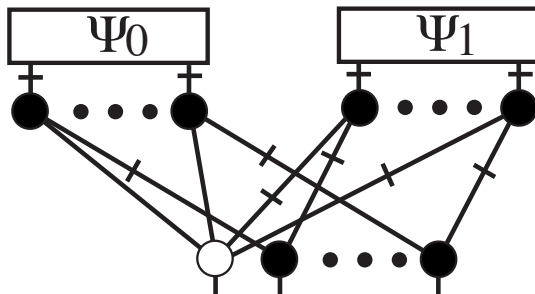
$$|\Psi\rangle_0 \otimes |\Psi\rangle_1 \longmapsto |0, \Psi_0\rangle + |1, \Phi_1\rangle$$

- ▶ But that's impossible, because it's not linear. However, this is SLOCC, so we only care up to a scalar
- ▶ We can do almost as good by just projecting out the unwanted state:

$$|\Psi\rangle \otimes |\Phi\rangle \longmapsto \langle p|\Phi\rangle |0, \Psi\rangle + \langle p|\Psi\rangle |1, \Phi\rangle$$

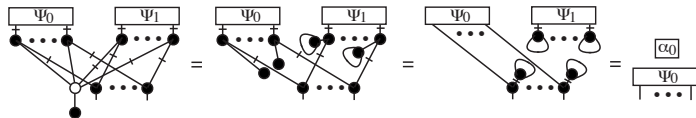
Inductive Classification

- ▶ We do this graphically like so:



Inductive Classification

- ▶ $(|\downarrow\rangle, |\uparrow\rangle)$ is a basis for $(\mathbb{C}^2)^*$, so we use them to project the control bit:



- ▶ ...and similarly, projecting by $|\uparrow\rangle$ yields $\alpha_1|\Psi_1\rangle$.

Conclusion

- ▶ Tripartite states define Frobenius algebras
- ▶ Since (up to SLOCC) there are two tripartite states over \mathbb{C}^2 , there are two types of commutative Frobenius algebras, special and anti-special
- ▶ In studying the interactions between a special Frobenius algebra and its associated anti-special FA, a rich algebraic theory crops up, which we are only beginning to understand
- ▶ We can already capture many behavioural properties of states, and provide an inductive technique for “building to order” n -partite states

Future Work

- ▶ We'd like to do this in arbitrary monoidal categories, not just FHilb . Axiomatize!
- ▶ The red-green theory can actually be encoded as a GHZ/W pair. The details need to be worked out.
- ▶ Questions of exhaustiveness, efficiency, and universality
- ▶ When and how can we go backwards from a vector to a graph?

Thanks!

This is joint work with Bob Coecke. Big chunks also due to Ross Duncan, Bill Edwards, and Simon Pedrix.

Questions?